QuickSort

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Like Merge Sort, QuickSort is a **Divide and Conquer algorithm**. It picks an element as pivot and partitions the given array around the picked pivot. There are many different versions of quickSort that pick pivot in different ways.

- 1. Always pick first element as pivot.
- 2. Always pick last element as pivot (implemented below)
- 3. Pick a random element as pivot.
- 4. Pick median as pivot.

The key process in quickSort is partition(). Target of partitions is, given an array and an element x of array as pivot, put x at its correct position in sorted array and put all smaller elements (smaller than x) before x, and put all greater elements (greater than x) after x. All this should be done in linear time.

Pseudo Code for recursive QuickSort function :

```
/* low --> Starting index, high --> Ending index */
quickSort(arr[], low, high)
{
    if (low < high)
    {
        /* pi is partitioning index, arr[pi] is now
        at right place */
        pi = partition(arr, low, high);
        quickSort(arr, low, pi - 1); // Before pi
        quickSort(arr, pi + 1, high); // After pi
    }
}</pre>
```

Partition Algorithm

There can be many ways to do partition, following pseudo code adopts the method given in CLRS book. The logic is simple, we start from the leftmost element and keep track of index of smaller (or equal to) elements as i. While traversing, if we find a smaller element, we swap current element with arr[i]. Otherwise we ignore current element.

Pseudo code for partition()

```
/* This function takes last element as pivot, places
   the pivot element at its correct position in sorted
    array, and places all smaller (smaller than pivot)
   to left of pivot and all greater elements to right
   of pivot */
partition (arr[], low, high)
{
    // pivot (Element to be placed at right position)
    pivot = arr[high];
    i = (low - 1) // Index of smaller element
    for (j = low; j <= high- 1; j++)</pre>
    {
        // If current element is smaller than the pivot
        if (arr[j] < pivot)</pre>
        {
            i++;
                    // increment index of smaller element
            swap arr[i] and arr[j]
        }
    }
    swap arr[i + 1] and arr[high])
    return (i + 1)
}
```

Example of Quick Sort:

44 33 11 55 77 90 40 60 99 22 88

Let 44 be the Pivot element and scanning done from right to left

Comparing **44** to the right-side elements, and if right-side elements are **smaller** than **44**, then swap it. As **22** is smaller than **44** so swap them.

22 33 11 55 77 90 40 60 99 44 88

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And t	hese s	sublists	are so	orted ur	der the	e same	proces:	s as abo	ove dor	ie.
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Swap	with 7	7:								
22	33	11	40	77	90	44	60	99	55	88
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22 3 Recur										